



# PARALLEL PROCESSING

## UNIT 4

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## QUIZ

On scan of  $N$  elements the amount of work is:

- $O(\log n)$
- $O(n)$
- $O(n \log n)$
- $O(n^2)$

On scan of  $N$  elements the number of steps is:

- $O(\log n)$
- $O(n)$
- $O(n \log n)$
- $O(n^2)$

# OUTLINES

- Compact
- Compact-like
- Segment scan
- Sorting
  - Odd Even Sort
  - Merge Sort
  - Radix Sort
  - Quick Sort



# COMPACT (FILTER)

4



# COMPACT MODEL

Input:            a      b      c      d      e      ...

Predicate:        T      F      T      F      T      ...

e.g.

“Is my index even?”

Output:           a      -      c      -      e      ... Sparse

                  a      c      e      ...            Dense



# QUIZ

When to use compact?

When the number of elements .....

- Small
- Large

When the operation on this elements are .....

- Expensive (complex)
- Cheap (Simple)



# COMPACT PARALLELIZATION

- How to compute the scatter address/index in parallel?

Input:                    a        b        c        d        e        ...

Predicate:                T        F        T        F        T        ...

e.g.

“Is my index even?”

Output:                    a                    -                    c                    -                    e                    ... Sparse  
                                  ↓                    ↙                    ↘  
                                  a                    c                    e                    ... Dense

# COMPACT PARALLELIZATION

- We can paraphrase the problem as following:

<b>Input:</b>	<b>T</b>	<b>F</b>	<b>F</b>	<b>T</b>	<b>T</b>	<b>F</b>
<b>Output:</b>	<b>0</b>	<b>-</b>	<b>-</b>	<b>1</b>	<b>2</b>	<b>-</b>

- And for computer we paraphrase it again as following:

<b>Input:</b>	<b>1</b>	<b>0</b>	<b>0</b>	<b>1</b>	<b>1</b>	<b>0</b>	<b>???</b>
<b>Output:</b>	<b>0</b>	<b>1</b>	<b>1</b>	<b>1</b>	<b>2</b>	<b>3</b>	

**Exclusive Scan**

# COMPACT PARALLEL ALGORITHM

1. Generate Predicate array:

```
predicate_array = predicate_function(input_array)
```

2. Generate Scan-in array (1s and 0s)

```
scan-in_array = convert(predicate_array)
```

3. Generate Scatter-addresses array :

```
addresses_array = exclusive_sum_scan(scan-in_array)
```

4. Scatter input elements to addresses

```
output = scatter(addresses_array, input_array)
```

} map

} scan

} scatter

# QUIZ

- Suppose we need to compact 1M element with the following predicate functions:
  - A: `isDivisibleBy_17( )`      Few elements
  - B: `isNotDivisibleBy_34( )`      Many elements

	A Faster	A = B	B Faster
Map		✓	
Scan		✓	
Scatter	✓		

# OUTLINES

- Compact
- Compact-like
- Segment scan
- Sorting
  - Odd Even Sort
  - Merge Sort
  - Radix Sort
  - Quick Sort

The background features a black field with vertical stripes on the left and right sides. The stripes consist of a thin black line, a wider textured grey line, and a thin orange line. Several orange circles of varying sizes are scattered across the page, including a large one on the left, a medium one in the center, and a small one on the right.

# COMPACT-LIKE

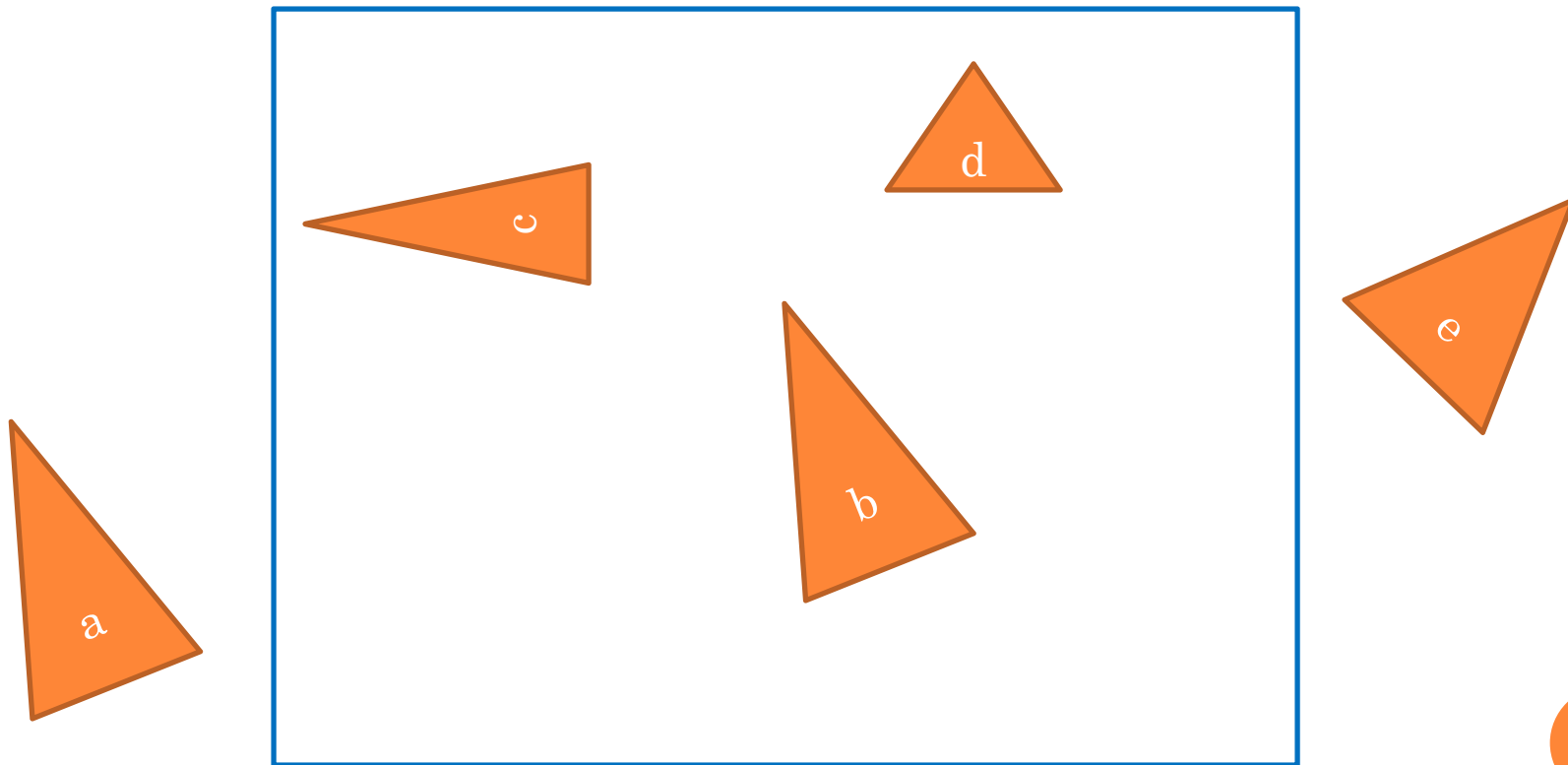
14

## RECAP COMPACT ALLOCATION

- Compact allocate 1 output for 1 (true) element and 0 output for 1 (false).
- Can we generalize?? (not only 1s)
  - The number of output can be computed dynamically for each input items.

# CLIPPING

- Suppose a set of triangles are sent as input to a computer graphics pipeline







# BAD SOLUTION

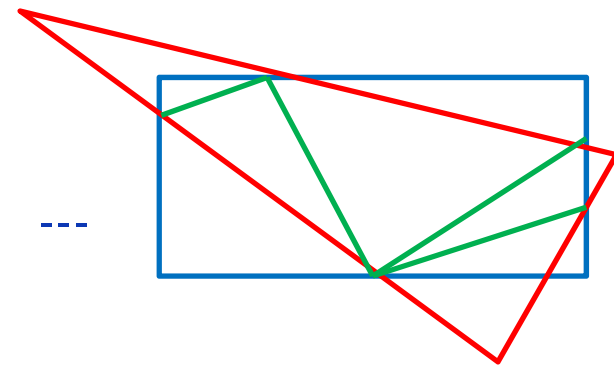
Input: 

a	c	e	d	f	b	g
---	---	---	---	---	---	---

- Allocate maximum possible space in intermediate array
  - 5 for each triangle in our case
  - **Intermediate:**

a	?	?	?	?	b	?
---	---	---	---	---	---	---

 ---
- Apply compact



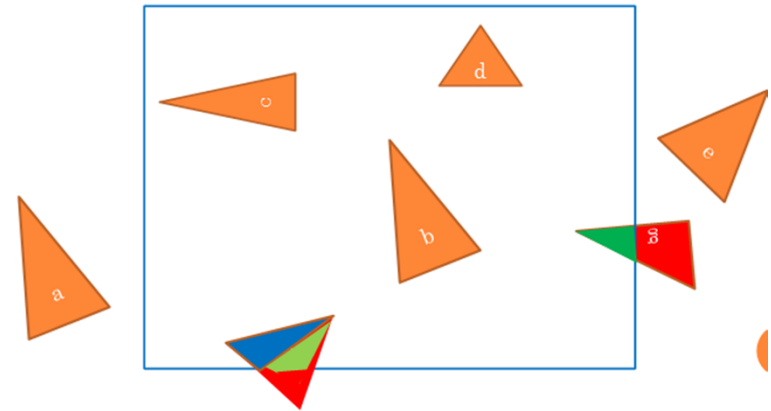
Disadvantages:

- Wasteful in space
- Scanning large intermediate array

# GOOD SOLUTION (GENERAL COMPACT)

Input: 

a	c	e	d	f	b	g
---	---	---	---	---	---	---



- Allocation requests per input element

request: 

0	1	0	1	2	1	1
---	---	---	---	---	---	---

- Apply scan

Addresses: 

0	0	1	1	2	4	5
---	---	---	---	---	---	---

- Allocate output array based with respect to max scan #, and then apply scatter

output: 

0	1	2	3	4	5
---	---	---	---	---	---

## OTHER APPLICATION OF SCAN

- Data Compression
- Collision Detection

# OUTLINES

- Compact
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The background features a black field with vertical stripes of varying widths and colors, including a wide grey stripe on the left and a thin white stripe. Several orange circles of different sizes are scattered across the page, with one large circle on the left and another on the bottom right.

# SEGMENTED SCAN

22

# SEGMENTED SCAN

- Many small Scans
  - Lunch each independently
  - Combine as segments
- Remember we back all segments in one big array to be processed by one kernel, instead of running a separate kernel over each segment to gain max benefits from GPU power.
- We use a separate array to indicate segments' heads

**Input:** 1 2 3 4 5 6 7 8  $\xrightarrow{\text{Scan}}$  0 1 3 6 10 15 21 28

**Heads:** 1 0 1 0 0 1 0 0

**Input:** 1 2 | 3 4 5 | 6 7 8  $\xrightarrow{\text{Segmented Scan}}$  0 1 | 0 3 7 | 0 6 13

# SEGMENTED SCAN APPLICATION

- “Sparse matrix \*Dense vector” multiplication (SpMv)
  - Sparse matrix: contains all elements includes a lot of zeros
  - Dense matrix: doesn't contain zeros
- Sparse matrix multiplication comes with a lot of unnecessary multiplication
- E.g. Google PageRank
  - a is a non-zero value indicate there is a link between webpages indicated by column index and webpage indicated by row index.

$$\begin{array}{c} \text{All web pages} \\ \left[ \begin{array}{ccc} a & \cdots & \\ \vdots & \ddots & \vdots \\ \vdots & \vdots & \vdots \end{array} \right] \end{array}$$





# REVIEW MATRIX MULTIPLICATION

$$\begin{bmatrix} 1 & 0 & 3 \\ 2 & 4 & -1 \\ 0 & 1 & 5 \end{bmatrix} \times \begin{bmatrix} 0 \\ 1 \\ 2 \end{bmatrix} =$$

$$\begin{bmatrix} (1 * 0) + (0 * 1) + (3 * 2) \\ (2 * 0) + (4 * 1) + (-1 * 2) \\ (0 * 0) + (1 * 1) + (5 * 2) \end{bmatrix} =$$

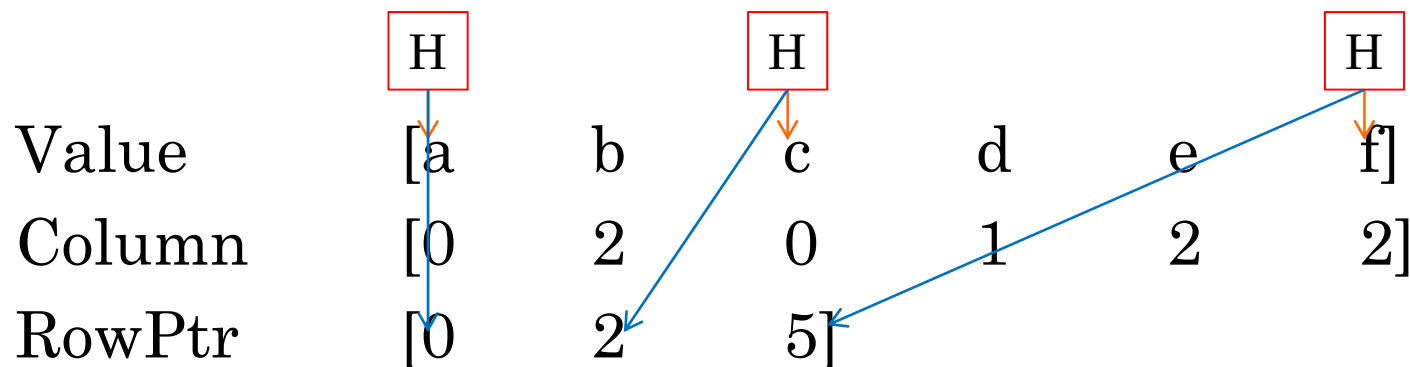
$$\begin{bmatrix} 6 \\ 2 \\ 11 \end{bmatrix}$$



# COMPRESSED SPARSE ROW

- We can represent sparse matrix in CPR format as following:

$$\begin{bmatrix} a & 0 & b \\ c & d & e \\ 0 & 0 & f \end{bmatrix}$$



# CPR MULTIPLICATION

○ We can represent sparse matrix in CPR format as following:

- Value: [a b c d e f]
  - Column: [0 2 0 1 2 2]
  - RowPtr: [0 2 5]
- $$\times \begin{bmatrix} x \\ y \\ z \end{bmatrix} \begin{matrix} 0 \\ 1 \\ 2 \end{matrix}$$

1. Create Segments with values and RowPtr

$$[a \ b \ c \ d \ e \ f] \xrightarrow[\text{RowPtr: } [0 \ 2 \ 5]]{\quad} [a \ b \mid c \ d \ e \mid f]$$

2. Gather vector values using column

$$\xrightarrow[\text{column: } [0 \ 2 \ 0 \ 1 \ 2 \ 2]]{\quad} [x \ z \ x \ y \ z \ z]$$

3. Pairwise multiply 1 . 2

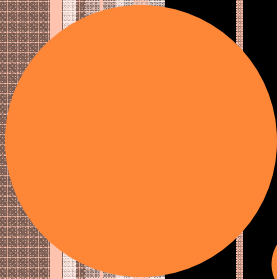
$$[a*x \ b*z \mid c*x \ d*y \ e*z \mid f*z]$$

4. Apply exclusive backward sum scan (at the head)

Can we apply reduce instead?? out(0) out(1) out(2)

# OUTLINES

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# SORTING

29

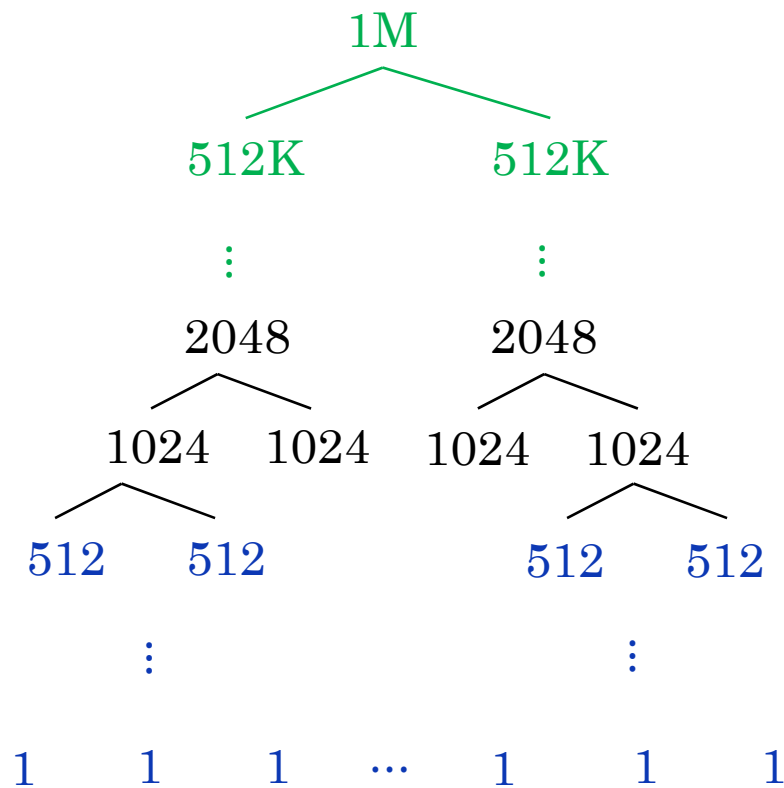


# ODD-EVEN (BRICK) SORT

- It's the parallel version of bubble sort

						Step?	Work?
	5	1	4	2	3		
1	5	2	4	3		<input type="checkbox"/>	<input type="checkbox"/>
	1	2	5	3	4	<input checked="" type="checkbox"/>	<input type="checkbox"/>
1	2	3	5	4		<input type="checkbox"/>	<input type="checkbox"/>
1	2	3	4	5		<input type="checkbox"/>	<input checked="" type="checkbox"/>

# MERGE SORT



## Stage 3:

- We have only 1 merge task
- 2 huge sorted lists
- Choose splitters (256<sup>th</sup> element) in each list. and sort them
- Use merge task in stage 2 to merge elements between splitters

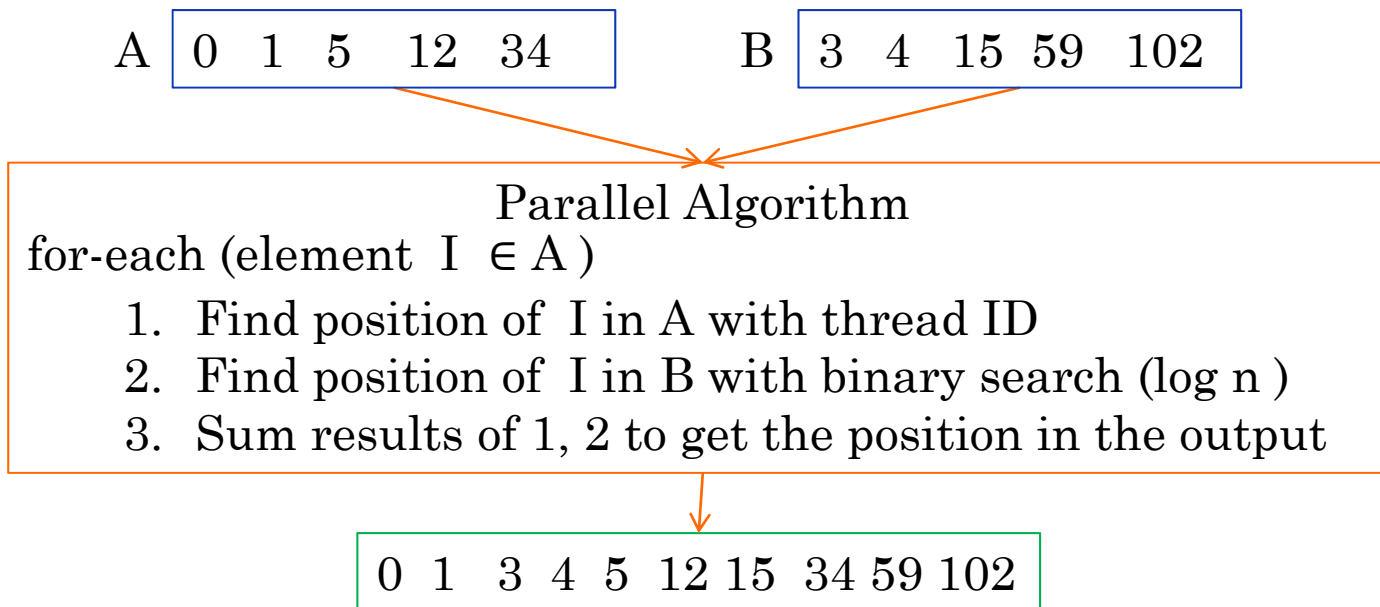
## Stage 2:

- Bunch of medium merge tasks
- 1 merge to 1 block

## Stage 1:

- Tons of small merge tasks
- 1 merge to 1 thread
- We can use shared memory
- Also we can use serial algorithm to sort small block of elements

# MERGE TASK

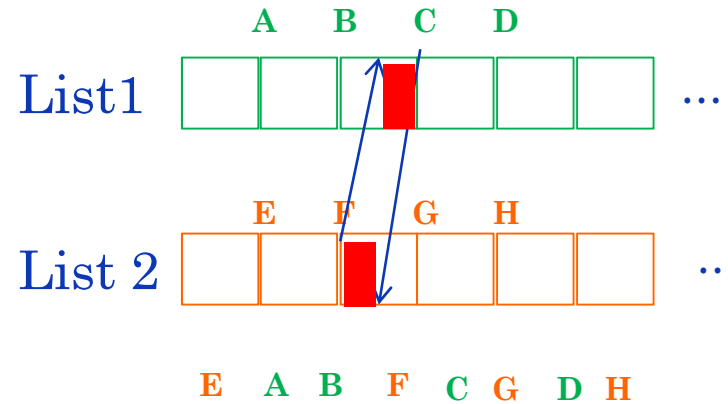




# MERGE TASK (CONT.)

- How to merge two huge lists in parallel?
  - No one can do all of this work

Find  $n^{\text{th}}$  (e.g.  $256^{\text{th}}$ ) splitter elements

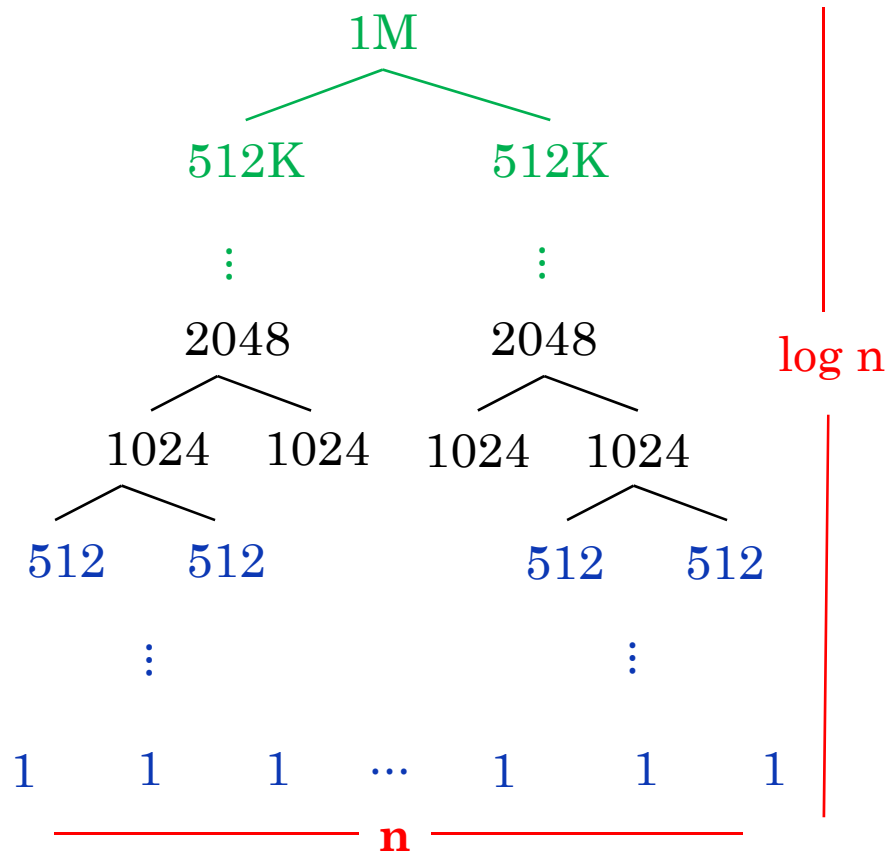


Merge splitters

Merge elements between every two consecutive splitters  
e.g. F and C:

1. Find F position in list 1
2. Find C position in list 2
3. Merge elements between these two positions in every list.

# MERGE ANALYSIS



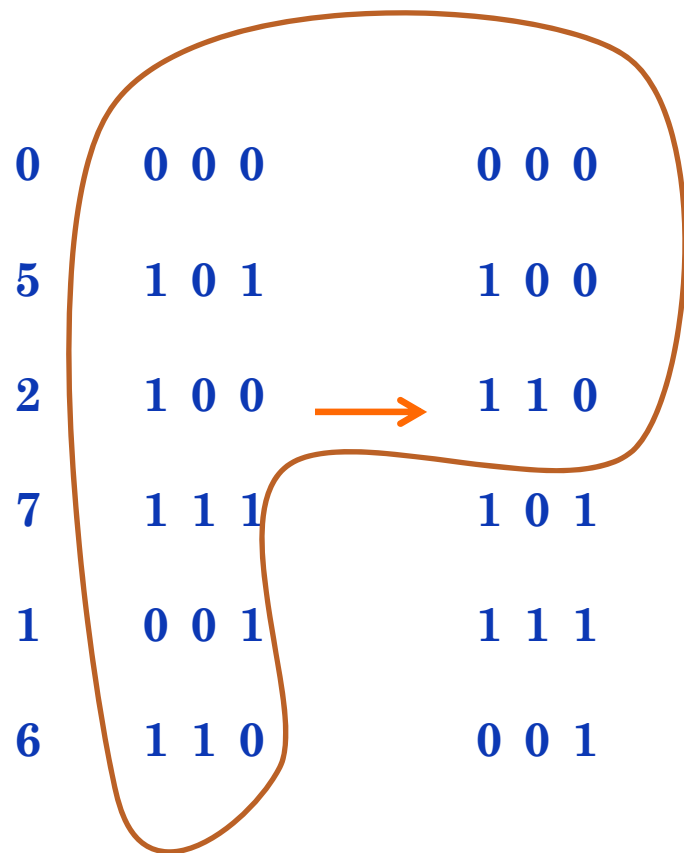
	Step?	Work?
$O(1)$	<input type="checkbox"/>	<input type="checkbox"/>
$O(n)$	<input type="checkbox"/>	<input type="checkbox"/>
$O(\log n)$	<input checked="" type="checkbox"/>	<input type="checkbox"/>
$O(n \log n)$	<input type="checkbox"/>	<input checked="" type="checkbox"/>
$O(n^2)$	<input type="checkbox"/>	<input type="checkbox"/>

# RADIX SORT

1. Start with LSB
2. Split input into 2 sets based on the current bit
3. Move to next MSB, and repeat 1

0	0 0 0	0 0 0	0 0 0	0 0 0	0
5	1 0 1	1 0 0	1 0 0	0 0 1	1
4	1 0 0	1 1 0	1 0 1	1 0 0	4
7	1 1 1	1 0 1	0 0 1	1 0 1	5
1	0 0 1	1 1 1	1 1 0	1 1 0	6
6	1 1 0	0 0 1	1 1 1	1 1 1	7

# RADIX SORT PENALIZATION



What is this Algorithm?

**Compact**

What is The Predicate?

**$(i \& 1) == 0$**



# RADIX SORT ANALYSIS

0	0 0 0	0 0 0
5	1 0 1	1 0 0
2	1 0 0	1 1 0
7	1 1 1	1 0 1
1	0 0 1	1 1 1
6	1 1 0	0 0 1

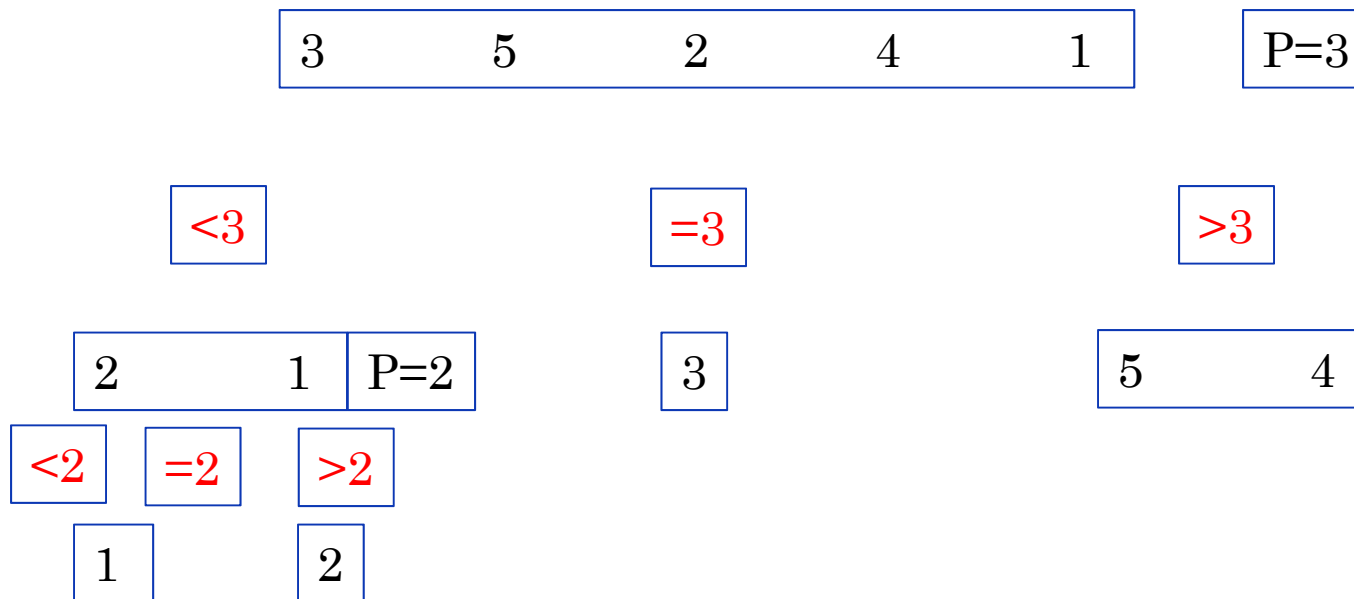
Work:  $O(kn)$   
linear

← k: #of bits      n: #of elements →

Steps:  $O(k)$

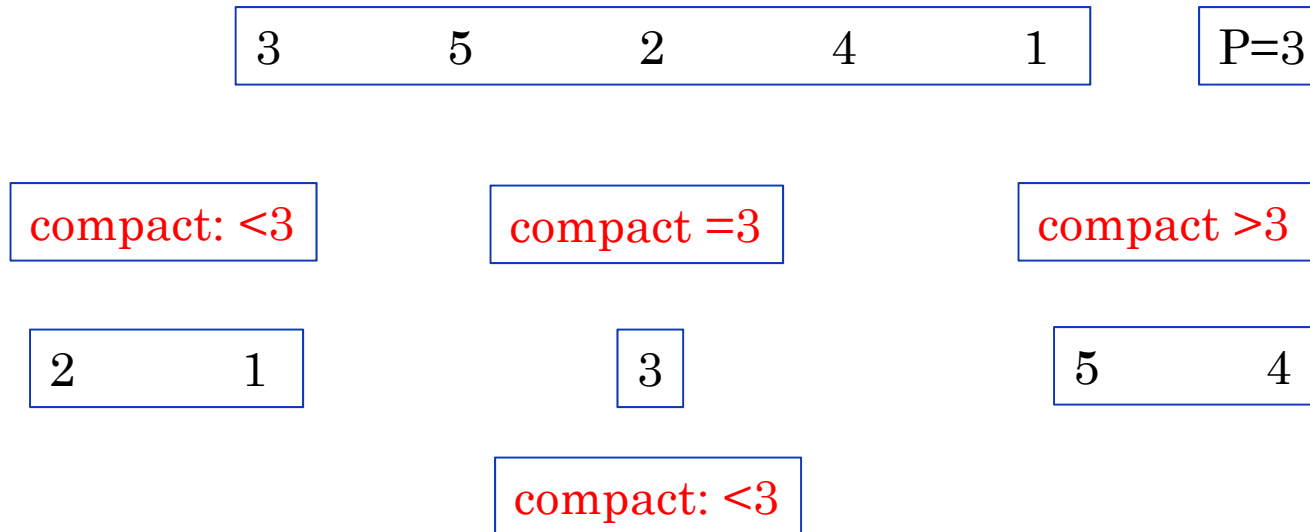
# QUICK SORT

- Choose Pivot element
- Compare all elements with Pivot
- Split into 3 Arrays  $<p$ ,  $=p$ ,  $>p$
- Recursion on each array



# QUICK SORT **PENALIZATION**

- Old GPUs doesn't support recursion



- Current GPUs support recursion

## NOTE

- All sort algorithms that we have studied are key value sorts where we usually depend on an integer key to sort.
- However if you have items with different data item (e.g. structure with many value).
  - Use a key or a pointer to this value to apply sorting



# OUTLINES

- Compact
- Compact-like
- Segment scan
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# RED EYE REMOVAL



- Stencil
- Sort
- Map